

# An SQP-type method for constrained and unconstrained nonlinear multiobjective optimization

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Recent Advances in Multi-Objective Optimization

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# Outline

- 1 Introduction
- 2 The algorithm
- 3 Implementation
- 4 Numerical results
- 5 Conclusions

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# Multiobjective constrained optimization

## Constrained multiobjective optimization problem

$$\min_{x \in \Omega} f(x) = (f_1(x), \dots, f_m(x))^T$$

with

$$\Omega = \{x \in [\ell, u] \subseteq \mathbb{R}^n : g_j(x) \leq 0, j = 1, \dots, p, \quad h_l(x) = 0, l = 1, \dots, q\}$$

- $\ell \in (\mathbb{R} \cup \{-\infty\})^n, u \in (\mathbb{R} \cup \{+\infty\})^n$ ;
- All objective functions are at least  $C^2$ ;
- All constraint functions are at least  $C^1$ ;
- We also allow **unconstrained** or box-constrained optimization ( $p, q = 0$  and/or  $\ell = \{-\infty\}^n, u = \{+\infty\}^n$ ).



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# Algorithm main lines

- Does **not aggregate** any of the objective functions
- Uses **SQP** based techniques for MOO
- Keeps a list of **nondominated** points
- Constraints **violations** are considered as additional objectives in the linesearch steps.
- Tries to capture the **whole Pareto front** from two algorithmic stages: search and refining

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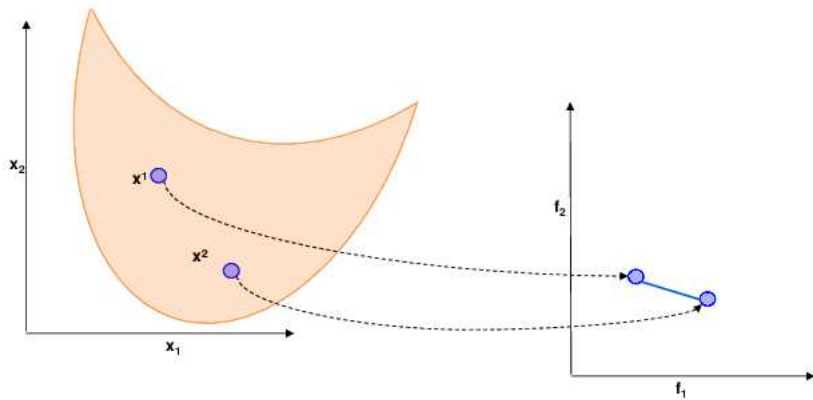
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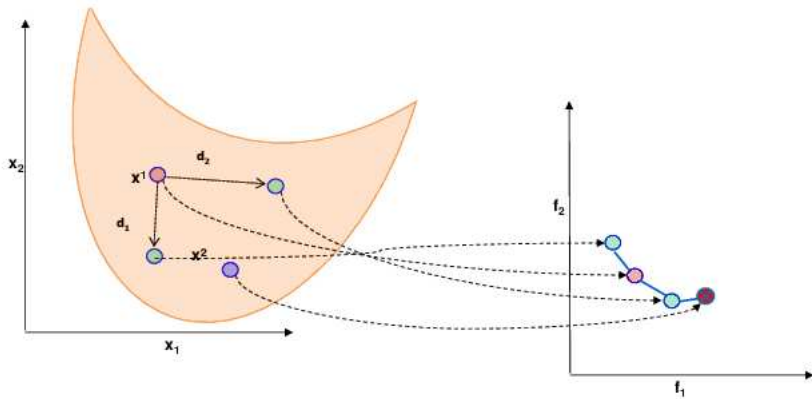
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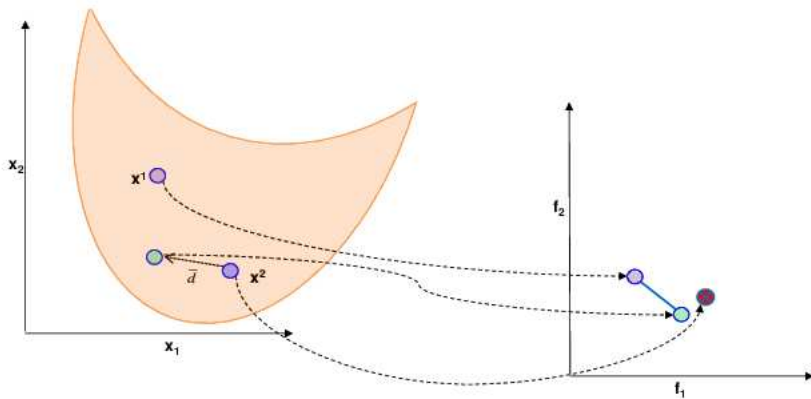
# Algorithm illustrated - setup



# Algorithm illustrated - spread



# Algorithm illustrated - refining



# Search direction computation

For each  $x_k$  in the list of nondominated points:

Spread ( $i = 1, \dots, m$ )

$$\begin{aligned} d_i \in \arg \min_{d \in \mathbb{R}^n} \quad & \nabla f_i(x_k)^T d + \frac{1}{2} d^T H_i d \\ \text{s.t.} \quad & g_j(x_k) + \nabla g_j(x_k)^T d \leq 0, \quad j = 1, \dots, p \\ & h_l(x_k) + \nabla h_l(x_k)^T d = 0, \quad l = 1, \dots, q \\ & \ell \leq x_k + d \leq u \end{aligned}$$

where  $H_i$  is a **positive definite** matrix.

$d_i$  is a descent direction for  $f_i$ .

Linesearch: test  $x_k + \alpha d_i$  ( $\alpha = 2^t$ ,  $t = 0, 1, 2, \dots$ ) for being nondominated.

# Search direction computation

For each  $x_k$  in the list of nondominated points:

Refining

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & \sum_{i=1}^m f_i(x) \\ \text{s.t.} \quad & f_i(x) \leq f_i(x_k), \quad i = 1, \dots, m \\ & g_j(x) \leq 0, \quad j = 1, \dots, p \\ & h_l(x) = 0, \quad l = 1, \dots, q \end{aligned}$$

Iterations of an SQP-type method for this problem are carried out, using  $x_k$  as a starting point.

# Some theoretical considerations

- From **spread stage** we obtain new points.
- The **spread stage** performs a finite number of iterations. (No asymptotics here!)
- The **refining stage** drives all the available nondominated points to Pareto criticality,
- by obtaining a new point that **decreases** or **maintains** all the objective function values.
- (Local) Pareto **criticality** can be verified based on the refining single-objective optimization problem.
- **Convergence theory** available for the proposed algorithm.



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# Implementation

- Implemented in MATLAB (fast prototyping, high performance)
- (Single-objective) **subproblems** are solved by quadprog and fmincon MATLAB solvers
- **Maximum** of 20 iterations on the spread stage
- We consider three possibilities for the  $H_i$  matrix:
  - $H_i = I_n$  in both stages
  - $H_i = \begin{pmatrix} I_{n-1} & 0 \\ 0 & 1 \end{pmatrix}$  in both stages ( $I_n$  p.d. matrix)
  - $H_i = I_n$  in the spread stage and  $H_i = \begin{pmatrix} I_{n-1} & 0 \\ 0 & 1 \end{pmatrix}$  in the refining stage
- Two (list) **initialization** strategies are implemented:



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  - $H_i = \nabla^2 \phi(x_i)$  in both stages (if not available,  $H_i = I_m$ )
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  - **Random initialization**: random initialization of the subproblems
  - **Warm-start initialization**: warm-start initialization of the subproblems

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- Two (list) **initialization** strategies are implemented:

• **1313a** – a line between  $\ell$  and  $w$  ( $x_i = \ell + i \frac{w-\ell}{2n\delta}$ ,  $i = 1, \dots, 2n\delta$ )

• **1313b** – a line between  $\ell$  and  $w$  ( $x_i = \ell + i \frac{w-\ell}{2n\delta}$ ,  $i = 1, \dots, 2n\delta$ )

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  - line – a line between  $\ell$  and  $u$  ( $x_i = \ell + i \frac{u-\ell}{2nS}$ ,  $i = 1, \dots, 2nS$ ).
  - rand – a uniform  $(\ell, u)$  random distribution.

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# Academic test set

- Problems from the academic literature.
- Problems have been coded in AMPL (and a MATLAB-AMPL interface was used). Exact derivatives are provided by AMPL.
- 67 bound constrained test problems (50 problems with  $m = 2$ , 17 problems with  $m = 3$ ),  $n$  varying between 2 and 30.
- 21 constrained test problems (12 problems with  $m = 2$ , 9 problems with  $m = 3$ ), 7 with nonlinear constraints, 9 with linear constraints, and 5 with both,  $n$  varying between 2 and 20.

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# Solvers

- We consider **six implementations** of the MOSQP solver: MOSQP ( $H = I$ , line), MOSQP ( $H = \nabla^2 f$ , line), MOSQP ( $H = (I, \nabla^2 f)$ , line), MOSQP ( $H = I$ , rand), MOSQP ( $H = \nabla^2 f$ , rand), MOSQP ( $H = (I, \nabla^2 f)$ , rand).
- MOSQP compared against **NSGA-II (C version)** and **MOScalar** (weighted-sum, equidistant weights).
- We report numerical results via performance and data profiles.

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• Performance profiles: *Purity*, *Spread*, *Gamma Metric*, *Spread Delta Metric*, and *Hypervolume metric*. (Small values  $\leftrightarrow$  better performance.)

• Data profiles: *best*, *mean*, *median*, *mean + std*, *best + std*, *median + std*.

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# Performance metrics

- *Spread-Gamma Metric*

For given solver and MOO problem, measures the largest gap in the approximate Pareto front.

- *Spread-Delta Metric*

For given solver and MOO problem, measures the uniformity of gaps in the approximate Pareto front.

- *Hypervolume*

For given solver and MOO problem, measures the volume of the space enclosed by the nondominated points and utopia point.

- *Purity*

For a given solver, a set of solvers, and a given MOO problem, measures  $1 / (\text{percentage of points computed that are not dominated by points from other solvers})$ .

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# Performance profiles

- Performance profiles.

Represent in one figure, for each solver  $s$ , the **cumulative distribution function**  $\rho_s$  for a given performance metric:

let  $\mathcal{P}$  bet the set of problems considered and  $r_{p,s}$  performance metric value of solver  $s$  on problem  $p$ . Then,

$$\rho_s(\tau) := |\{p \in \mathcal{P} : r_{p,s} \leq \tau\}|/|\mathcal{P}|.$$

- Solvers with larger  $\lim_{\tau \rightarrow \infty} \rho_s(\tau)$  are more robust. If  $\rho_s(\tau) = 1$  then solver  $s$  solves all given problems.

# Performance profiles

- Performance profiles.

Represent in one figure, for each solver  $s$ , the **cumulative distribution function**  $\rho_s$  for a given performance metric:

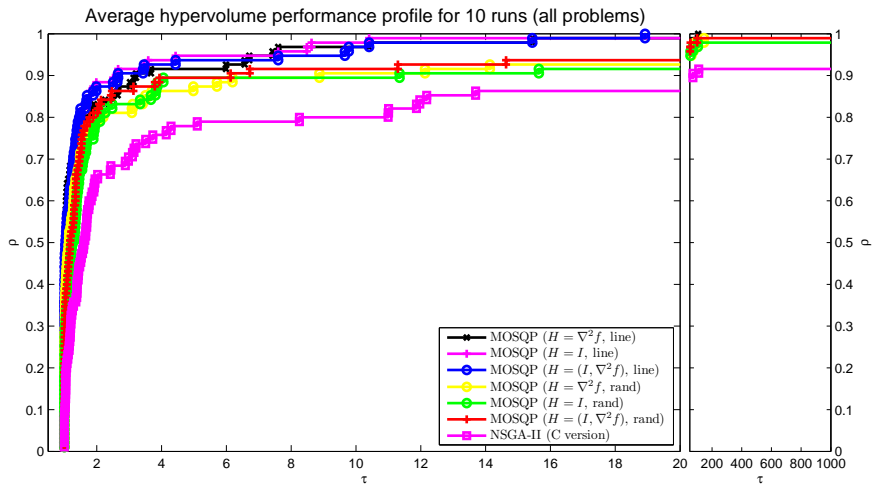
let  $\mathcal{P}$  bet the set of problems considered and  $r_{p,s}$  performance metric value of solver  $s$  on problem  $p$ . Then,

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- Solvers with **larger**  $\lim_{\tau \rightarrow \infty} \rho_s(\tau)$  are more robust. If  $\rho_s(\tau) = 1$  then solver  $s$  solves all given problems.

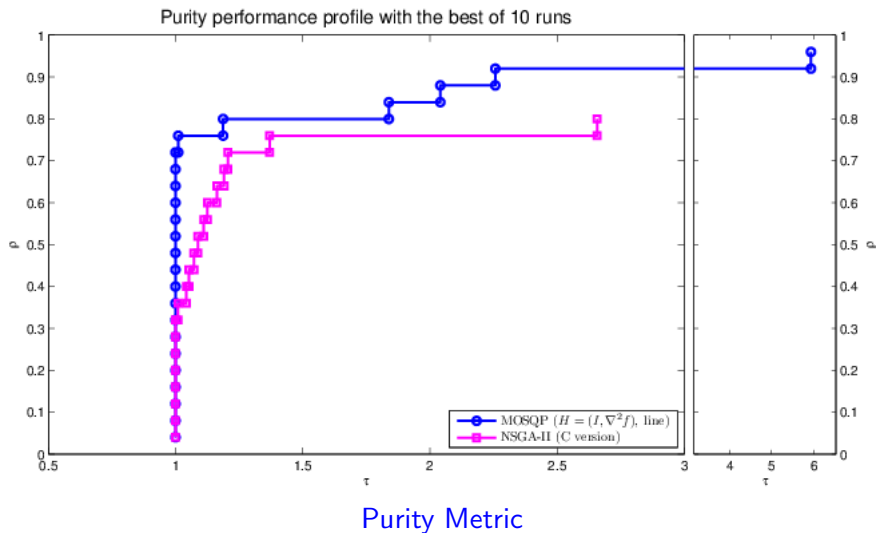


# Constrained test set (Hypervolume)

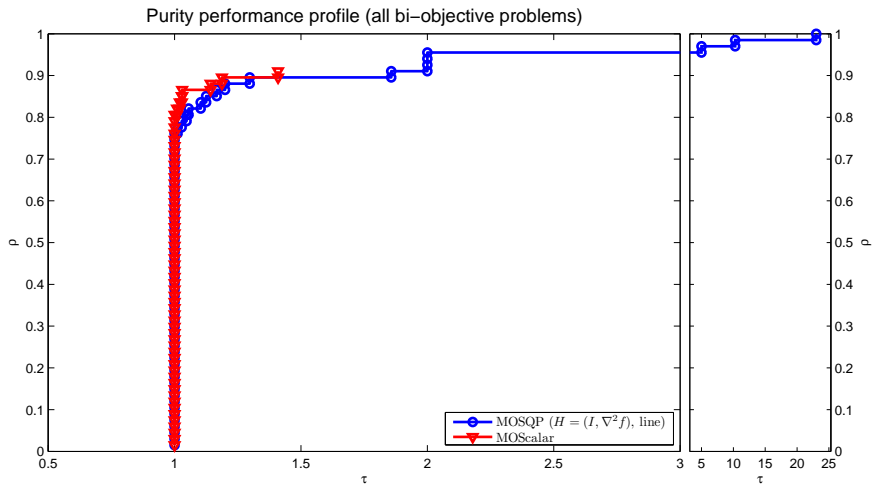


Hypervolume Metric

# Constrained test set (Purity)

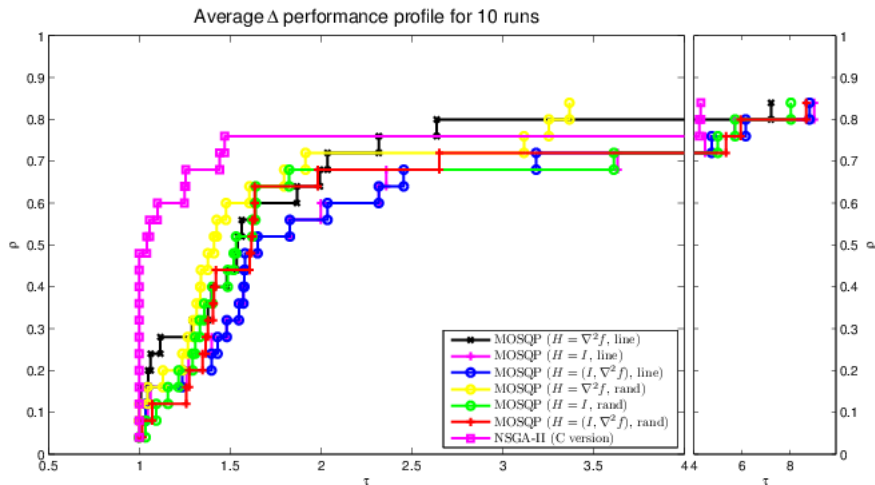


# Constrained test set (Purity)



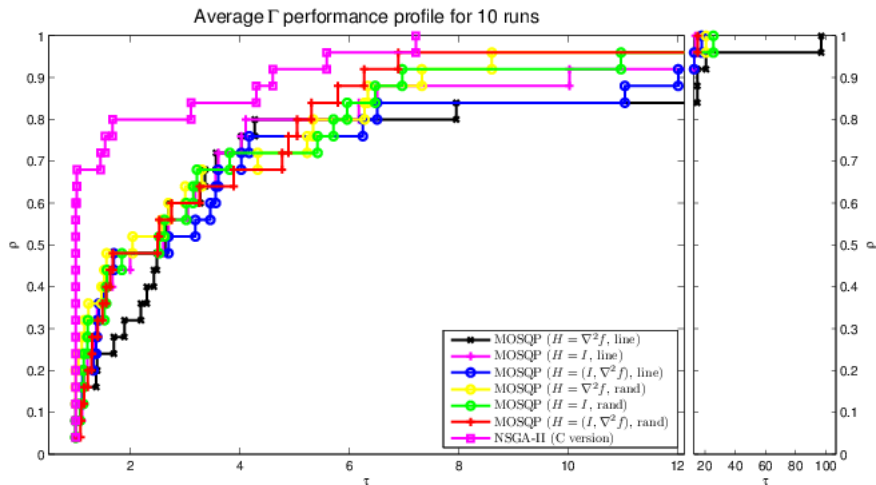
Purity Metric

# Constrained test set (Spread)



Delta Metric (uniformity of gaps in the Pareto front)

# Constrained test set (Spread)



Gamma Metric (largest gap in the Pareto front)

# Data profiles

Indicate how likely an algorithm is to ‘solve’ a problem, given some computational budget.

Let  $N_{p,s}$  be the number of function and gradient evaluations required for solver  $s$  to solve problem  $p$ :

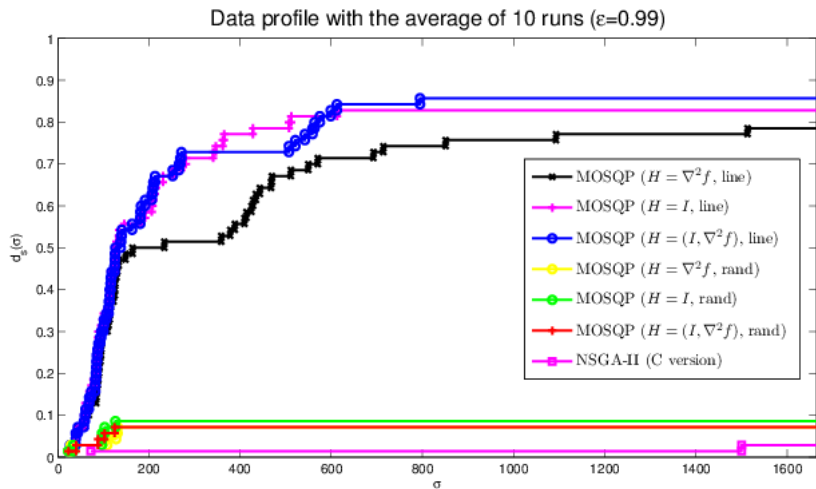
$$N_{p,s,f} := \sum_i \left( \#f_i + n\#\nabla f_i + \frac{(n+1)(n+2)}{2} \#\nabla^2 f_i \right),$$

$$N_{p,s} := N_{p,s,f} + N_{p,s,g} + N_{p,s,h}.$$

Consider the metric

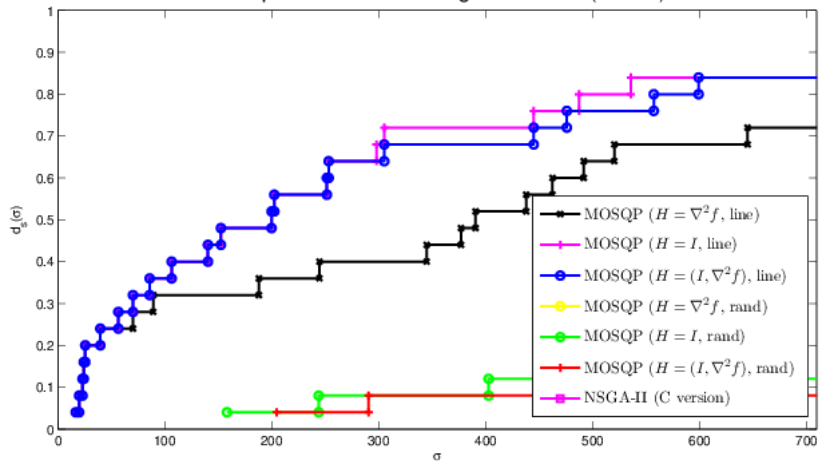
$$d_s(\sigma) = |\{p \in \mathcal{P} : N_{p,s} \leq \sigma\}| / |\mathcal{P}|.$$

# Bound constrained test set



# Constrained test set

Data profile with the average of 10 runs ( $\varepsilon=0.99$ )

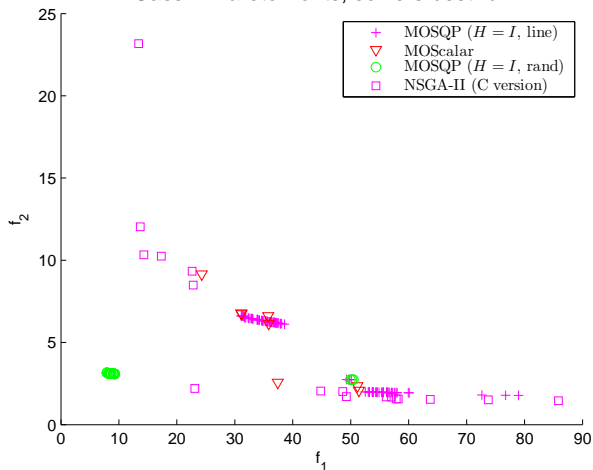




# Space Engineering: Earth-Jupiter Mission

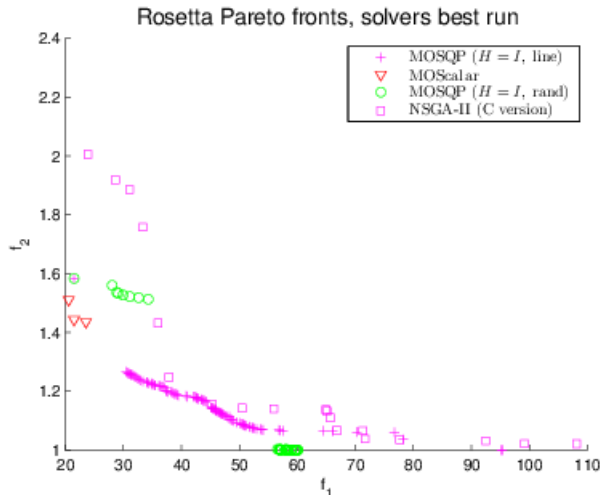
$f_1 = \Delta V \sim$  fuel use; and  $f_2 =$  total travel time.

Cassini Pareto fronts, solvers best run



# Space Engineering: Rosetta bi-objective problem

$f_1 = \Delta V \sim$  fuel use; and  $f_2 =$  total travel time.



# Outline

- 1 Introduction
- 2 The algorithm
- 3 Implementation
- 4 Numerical results
- 5 Conclusions**

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- We propose a method for constrained multi-objective optimization based on SQP, (**MOSQP**).
- Convergence proof establishes fast local convergence to Pareto front.
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- **Numerical results** confirm the solver competitiveness and robustness.

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# Thanks – Support

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